Railway Optimization

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- 1 Railway Planning
- 2 Line Planning Problem
- 3 Train Timetabling Problem





- 2 Line Planning Problem
- 3 Train Timetabling Problem





. . . back in the Greece

c. 600 BC - A basic form of the railway, the rutway, - existed in ancient Greek and Roman times, the most important being the ship trackway Diolkos across the lsthmus of Corinth. Measuring between 6 and 8.5 km, remaining in regular and frequent service for at least 650 years, and being open to all on payment, it constituted even a public railway, a concept which according to Lewis did not recur until around 1800. The Diolkos was reportedly used until at least the middle of the 1st century AD, after which no more written references appear.

• Timeline of railway history; From Wikipedia, the free encyclopedia





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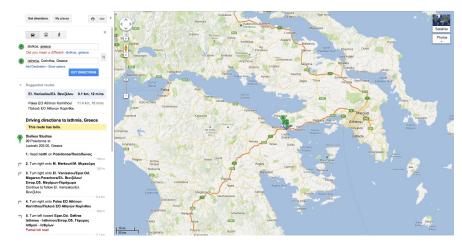
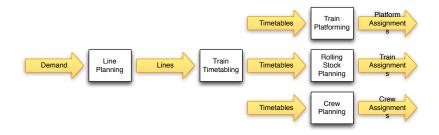


Figure : Source – Google Maps





Railway Planning Horizon







Day 0



Figure : Canadian Soldiers Building a Light Railway, [ca.1918]¹





- source: Archives of Ontario, Canadian Expeditionary Force Albums

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References



A. Caprara, L. G. Kroon, M. Monaci, M. Peeters, and P. Toth, *Passenger railway optimization*, Handbooks in Operations Research and Management Science (C. Barnhart and G. Laporte, eds.), vol. 14, Elsevier, 2007, pp. 129–187 (English).



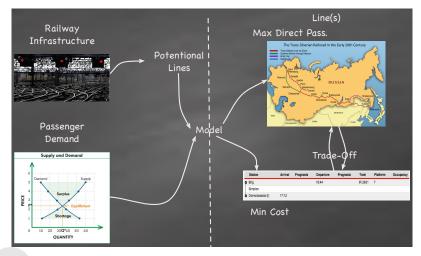


- 2 Line Planning Problem
- 3 Train Timetabling Problem





Line Planning Problem (LPP)







LPP Model I

G = (V, E)	-	undirected graph G representing the railway network
$v \in V$	-	set of stations
$e \in E$	-	set of edges representing the tracks between stations
$p \in P$	_	set of unordered pairs of stations $(p = (p_1, p_2))$ with positive demand
dp	-	number of passengers, that want to travel between stations p_1 and p_2
É _p	-	set of edges on the shortest path between stations p_1 and p_2
$d_e = \sum_{p:e \in E_p} d_p$	-	the total number of passengers, that want to travel along edge e
$I \in L$	-	set of potential lines (assumed to be known a priori)
E	-	set of edges of the line /
$f \in F$	-	set of potential frequencies
$c \in C$	-	set of available capacities
$i \in I$	-	set of indices representing combination of assigned capacity c
		and frequency f to a line l
ki	-	operational cost for a combination i (<i>e.g.</i> train driver, conductor(s), carriage kilometers
(1.4	£	

 $x_i = \begin{cases} 1 & \text{if and only if line } l_i \text{ is to be operated with a frequency } f_i \text{ and capacity } c_i, \\ 0 & \text{otherwise.} \end{cases}$

$$d_{lp}$$
 – number of direct passengers traveling on line l between the pair of stations p





LPP Model II

$$\begin{array}{ll} \max & w_1 \cdot \sum_{l \in L} \sum_{p \in P} d_{lp} - w_2 \cdot \sum_{i \in I} k_i \cdot x_i & (1) \\ \text{s.t.} & \sum_{i \in I \cdot l_i = I} x_i \leq 1, & \forall I \in L, \\ & (2) \\ & & (2) \\ \sum_{i \in I \cdot e \in E_{l_i}} f_i \cdot c_i \cdot x_i \geq d_e, & \forall e \in E, \\ & & (3) \\ & \sum_{p \in P \cdot e \in E_p} d_{lp} \leq f_i \cdot c_i \cdot x_i, & \forall I \in L, \forall e \in E, \\ & & (4) \\ & \sum_{l \in L \cdot E_p \subset E_l} d_{lp} \leq d_p, & \forall p \in P, \\ & & (5) \\ & x_i \in \{0, 1\}, & \forall i \in I, \\ & (6) \\ & d_{lp} \geq 0, & \forall I \in L, \forall p \in P. \\ & (7) \end{array}$$





2 Line Planning Problem

3 Train Timetabling Problem

- Non-Cyclic
- Cyclic





Train Timetabling Problem (TTP)

Non-Cyclic

- departs differently over the time horizon
- prior knowledge of the timetable needed
- lower cost

Cyclic

- departs at every cycle
- good for "unplanned" user
- higher cost





2 Line Planning Problem

3 Train Timetabling Problem

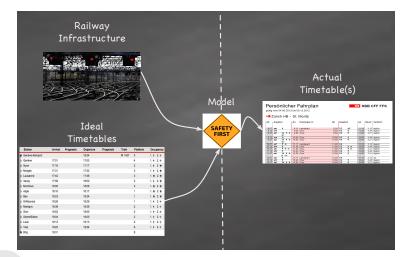
Non-Cyclic

Cyclic





Non-Cyclic TTP







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Non-Cyclic TTP Model I

G = (V, A)	_	directed acyclic multigraph G
$v \in V$	-	set of nodes
$a \in A$	-	set of arcs
$t \in T$	-	set of trains
$a \in A^t$	-	subset of arcs used by train t
σ	-	source node
au	-	sink node
pa	-	profit of arc a
$\delta_{t}^{+}(\mathbf{v})$	-	set of arcs in A^t leaving the node v
$\delta_t^-(\mathbf{v})$	-	set of arcs in A^t entering the node v
\mathbb{C}	-	family of maximal subsets C of pairwise incompatible arcs

 $x_a = \begin{cases} 1 & \text{if and only if the path in the solution associated with train } t \text{ contains arc } a, \\ 0 & \text{otherwise.} \end{cases}$





Non-Cyclic TPP Model II

$$\begin{array}{ll} \max & \sum_{t \in \mathcal{T}} \sum_{a \in A^{t}} p_{a} \cdot x_{a} & (8) \\ \text{s.t.} & \sum_{a \in \delta^{+}_{t}(\sigma)} x_{a} \leq 1, & \forall t \in \mathcal{T}, \\ & & (9) \\ & \sum_{a \in \delta^{-}_{t}(v)} x_{a} = \sum_{a \in \delta^{+}_{t}(v)} x_{a}, & \forall t \in \mathcal{T}, \forall v \in V \setminus \{\sigma, \tau\} \\ & & (10) \\ & \sum_{a \in C} x_{a} \leq 1, & \forall C \in \mathbb{C}, \\ & & (11) \\ & x_{a} \in \{0, 1\}, & \forall a \in A. \\ & & (12) \end{array}$$





2 Line Planning Problem

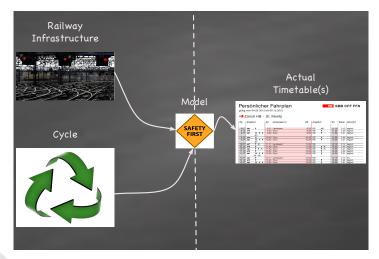
3 Train Timetabling Problem ■ Non-Cyclic

Cyclic





Cyclic TTP







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Cyclic TTP Model I

$G = (N, A \cup A^s)$	-	graph G representing the railway network
$n, m \in N$	-	set of nodes
$a \in A$	-	set of regular tracks $a = (n, m)$
$a \in A^s$	-	set of single tracks $a = (n, m) = (m, n)$
$t \in T$	-	set of trains
		set of nodes visited by train t
$a\in A^t\subseteq A\cup A^s$	-	set of tracks used by train t
$(t, t') \in T_a$	-	set of all pairs of trains (t, t') , that travel along the track <i>a</i> in the same direction, where t' is the faster train
$(t, t') \in T_{2}^{s}$	-	set of all pairs of trains (t, t') , that travel along the single
		track a in the opposite direction, where t departs from n
		and t' departs from m
$t \in F_n^d, F_n^a$	-	set of all trains t, that have a fixed departure (arrival) at node n
$(t, t') \in S_n$	-	set of all train pairs $(t, t'), t < t'$, for which the departure
		times are to be synchronized at node n
$(t, t') \in C_n$	-	set of all train pairs $(t, t'), t < t'$, for which turn-around or
		connection constraint is required from train t to train t' at node n
Ь	-	cycle of the timetable
h	-	general headway upon departure and arrival at every node
r ^t a	-	time it takes to the train t to traverse the arc a
$\begin{bmatrix} r_a^t \\ \underline{d_n^t}, \overline{d_n^t} \end{bmatrix}$ $\begin{bmatrix} \underline{f_n^t}, \overline{f_n^t} \end{bmatrix}$	-	dwell time window of the train t at the node n
$\left[\frac{f_n^t}{n}, \overline{f_n^t}\right]$	-	fixed arrival/departure window of the train t at the node n ,
		in the case of completely fixed arrival/departure $f_n^t = \overline{f_n^t}$
$\left[\frac{s_n^{tt'}}{s_n^{tt'}}, \overline{s_n^{tt'}}\right]$	-	time window for the synchronization of trains t and t] at node n
$\left[\frac{c_n^{tt'}}{c_n^{tt'}}, \overline{c_n^{tt'}}\right]$	-	time window for the connection or turn around constraint
		between trains t and t' at node n

- a_n^t arrival time of train t at node n
- d_n^t departure time of train t from node n





Cyclic TPP Model II

max	F(a, d)		(13)
s.t.	$a_m^t - d_n^t = r_a^t$	$\mod b$,	$\forall t \in T, a \in A^t$
			(14)
	$d_n^t - a_n^t \in \left[\underline{d}_n^t, \overline{d}_n^t\right]$	$\mod b$,	$\forall t \in T, \forall n \in N^t$,
			(15)
	$d_n^{t'} - d_n^t \in \left[\underline{s_n^{tt'}}, \overline{s_n^{tt'}}\right]$	mod b,	$\forall n \in N, \forall (t, t') \in S_n,$
			(16)
	$d_n^{t'} - a_n^t \in \left[c_n^{tt'}, \overline{c_n^{tt'}}\right]$	mod b,	$\forall n \in N, \forall (t, t') \in C_n,$
	,		(17)
	$d_n^{t'} - d_n^t \in \left[r_a^t - r_a^{t'} + h, b - h\right]$	mod b,	$\forall a \in A, \forall (t, t') \in T_a,$
			(18)
	$a_n^{t'} - d_n^t \in \left[r_a^t + r_a^{t'} + h, b - h\right]$	mod b,	$\forall a \in A^s, \forall (t, t') \in T_a,$
	L' ' .		(19)
	$d_n^t \in \left[f_n^t, \overline{f_n^t}\right]$,		$\forall n \in N, \forall t \in F_n^d$,
	L— 1		(20)
	$a_n^t \in \left[f_n^t, \overline{f_n^t}\right]$,		$\forall n \in N, \forall t \in F_n^a$,
			(21)
	$a_{n}^{t}, d_{n}^{t} \in \{0, b-1\},\$		$\forall t \in T, \forall n \in N_t.$
			(22)







Thank you for your attention.



Exam - 21.06.2013 (Friday)

Place GC A1 416

Structure

- 20 min presentation
- 20 min Q&A

Schedule

8:15 - Group 1 9:00 - Group 2 9:45 - Group 3 10:45 - Group 4 11:30 - Group 5 12:15 - Group 6



